Open Source in Industry:
All about mainline real-time Linux

Technical Heidelberg OSADL Talks, April 29, 2020, Online Session 1

History and functionality of real-time Linux
Functionalities of PREEMPT_RT
Linux Foundation RTL Collaborative Project
Live real-time demonstrator
Some information on today’s sessions

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• You may join an online discussion on all topics of today at 4 pm
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  ➢ Meeting name OSADLTechnicalHOT
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(We will show this slide again at the end of this session)
Quote of the year 2004

“It is *impossible* to turn a General Purpose Operating System Kernel into a RealTime Operating System Kernel.”
What is a real-time extension?

Hardware abstraction layer (nanokernel)

Real-time kernel infrastructure

Real-time application in kernel

Non-real-time kernel infrastructure

Real-time application #1

Real-time application #2

Real-time application #N

Non-real-time application #1

Non-real-time application #2

Non-real-time application #N

I/O

Interrupts

Hardware

Hardware abstraction layer (nanokernel)

Kernel

Userspace

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Open Source Automation Development Lab (OSADL), Heidelberg
What is a real time operating system?

Hardware abstraction layer

Real-time kernel infrastructure

Hardware

Interrupts

I/O

POSIX API

Any application #1

Any application #2

Any application #N

Kernel

Userspace
Terminology

- RTLinux (Dual-kernel approach) uses its own nanokernel RTAI, Xenomai, RTCore etc.

- CONFIG_PREEMPT_RT (Single-kernel approach) Linux mainline realtime, “Linux RTOS”
Goals

• Fully preemptive kernel
• Realtime guarantees suitable for the vast majority of applications
• POSIX compliance (single API):
  `sched_setscheduler()`
  `sched_setaffinity()`
  and friends
Driving forces

• Symmetric multi-processing
• (Bug fixing of race conditions)
• Audio recording and audio processing
• Video recording and video processing
• Reliable time-stamps in financial transactions
• Machine industry and embedded systems
People behind real-time Linux

- Doug Niehaus, University of Kansas
- Thomas Gleixner, Linutronix
- Ingo Molnàr, Red Hat
- Peter Zijlstra, Red Hat
- Paul E. McKenney, IBM
- Steven Rostedt, Red Hat
- many other
History

• Autumn 2004 MontaVista, Timesys, Lynuxworks post realtime related patch fragments

• Ingo Molnar re-implements parts from scratch and posts the real-time preemption patch

• A core team forms

• Kernel Summit 2006 in Ottawa accepts a plan to merge all components into mainline over time
Kernel-Summit, Ottawa, August 2006
"Controlling a laser with Linux is crazy, but everyone in this room is crazy in his own way. So if you want to use Linux to control an industrial welding laser, I have no problem with your using PREEMPT_RT."

Linus Torvalds
Main components

- Deterministic Scheduler
- Preemption Support
- PI Mutexes
- High-Resolution Timer
- Preemptive Read-Copy Update
- IRQ Threads (selected, forced)
- Raw Spinlock Annotation
- Preemptive Memory Management
- Full Realtime Preemption Support
<table>
<thead>
<tr>
<th>Feature</th>
<th>x86</th>
<th>x86/64</th>
<th>powerpc</th>
<th>arm</th>
<th>mips</th>
<th>68knommu</th>
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<td>Preemptive Read-Copy Update</td>
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<tr>
<td>Raw Spinlock Annotation</td>
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<tr>
<td>Full Realtime Preemption</td>
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<td>✔️</td>
<td>✔️</td>
<td>✔️</td>
<td>✔️ (3)</td>
</tr>
</tbody>
</table>

- Available in mainline Linux
- Available when Realtime-Preempt patches applied

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Configuration

Preemption Mode
Use the arrow keys to navigate this window or press the hotkey of the item you wish to select followed by the <SPACE BAR>. Press <?> for additional information about this option.

- No Forced Preemption (Server)
- Voluntary Kernel Preemption (Desktop)
- Preemptible Kernel (Low-Latency Desktop)
- Complete Preemption (Real-Time)
Latency measurement

\[ \text{cyclictest} \]
\[ \text{Mess.-Thread} \]
\[ T_1 \]
\[ T_2 \]
\[ T_N \]
Latency measurement

- Run on CPU #0
- Priority 99
- High-resolution timer
- Start one thread
- Interval 100 μs

```
# cyclicstest -n -a0 -t1 -p99 -i100
```

```
T: 0 (26623) P:99 I:100 C:2244063 Min:  2 Act:  5 Avg:  4 Max:   27
```

- Process ID
- Cycles
- Worst-case latency in μs
Latency measurement

Run on CPU #0

Priority 99, 98, 97, ... 88

High-resolution timer

Start 12 threads

Interval 100 μs

No delay

# cyclicetest -n -a0 -t12 -p99 -i100 -d0

T: 0 ( 2910) P:99 I:100 C:3217008 Min: 2 Act: 6 Avg: 4 Max: 32
T: 1 ( 2911) P:98 I:100 C:3217008 Min: 1 Act: 4 Avg: 3 Max: 59
T: 2 ( 2912) P:97 I:100 C:3217007 Min: 2 Act: 4 Avg: 3 Max: 47
T: 3 ( 2913) P:96 I:100 C:3217007 Min: 2 Act: 11 Avg: 3 Max: 53
T: 5 ( 2915) P:94 I:100 C:3217007 Min: 3 Act: 9 Avg: 7 Max: 89
T: 6 ( 2916) P:93 I:100 C:3217007 Min: 2 Act: 5 Avg: 4 Max: 85
T: 7 ( 2917) P:92 I:100 C:3217006 Min: 2 Act: 10 Avg: 5 Max: 119
T: 8 ( 2918) P:91 I:100 C:3217006 Min: 2 Act: 13 Avg: 9 Max: 148
T: 9 ( 2919) P:90 I:100 C:3217007 Min: 1 Act: 4 Avg: 4 Max: 178
T:10 ( 2920) P:89 I:100 C:3217006 Min: 1 Act: 4 Avg: 3 Max: 1413
T:11 ( 2921) P:88 I:100 C:3217006 Min: 3 Act: 7 Avg: 10 Max: 27331
Latency measurement

Run on all CPUs
High-resolution timer
Start 12 threads
Priority 99 of all threads
Interval 100 μs
No delay

# cyclictest -S -p99 -i100 -d0

T: 0 (15350) P:99 I:100 C:3839755 Min: 2 Act:  6 Avg:  3 Max:  24
T: 3 (15353) P:99 I:100 C:3839755 Min: 2 Act:  5 Avg:  4 Max:  24
T: 6 (15356) P:99 I:100 C:3839755 Min: 2 Act:  5 Avg:  3 Max:  17
T: 8 (15358) P:99 I:100 C:3839755 Min: 2 Act:  5 Avg:  4 Max:  22
T: 9 (15359) P:99 I:100 C:3839754 Min: 2 Act:  5 Avg:  4 Max:  22
T:10 (15360) P:99 I:100 C:3839755 Min: 2 Act:  5 Avg:  4 Max:  42
T:11 (15361) P:99 I:100 C:3839755 Min: 2 Act:  5 Avg:  5 Max:  34
Conclusion (1)

“It is possible to turn a General Purpose Operating System Kernel into a Realtime Operating System Kernel.”

Linux is a Real-Time Operating System (RTOS) now.
Conclusion (2)

- The real-time capabilities of the mainline Linux kernel are comparable to those that can be achieved using a traditional RTOS.
- The Linux real-time concept allows to switch to real-time without any major changes to existing drivers or user-space code.
Conclusion (3)

- For the time being, however, the mainline kernel does not contain all needed components, but a patch (about 750 kByte) still needs to be applied.
- The maintenance and the final merging of the patches to mainline is funded by a Linux Foundation project. It is expected that an important progress towards complete mainline integration will happen in kernel 5.8.
Real-time demonstrator using square wave signal
OSADL Parport adapter (https://www.osadl.org/?id=1575)
OSADL Parport adapter
Shell interface

```
echo 0 .. 255 >/dev/setparport set output byte
echo 256 .. 511 >/dev/setparport boolean "or" action with output byte
echo 512 >/dev/setparport clear all output bits
echo 513 >/dev/setparport set all output bits
echo 514 >/dev/setparport invert output bits
echo 515 >/dev/setparport increment output bits
echo 516 >/dev/setparport decrement output bits
echo 517 >/dev/setparport copy status register to output bits
echo 518 >/dev/setparport copy jiffies LSBs >> 10 to output bits
```
OSADL Parport adapter
NMI interface

The 4 input bits of the parallel port interface can be used for polling buttons when NMI triggers arrive. This makes it possible to execute particular actions such as SysReq commands in order to diagnose a crashed system. In addition NMIs can be used to generate LED display codes such as 515 to increment the displayed number:

```
modprobe setparport actions=yes nmicode=515
```
or
```
echo Y > /sys/module/setparport/parameters/actions
echo 515 > /sys/module/setparport/parameters/nmicode
```
OSADL Parport adapter
C interface

```c
char *device = "/dev/setparport";
int port = open(device, O_RDWR);
int count = 1;

char *value = "0";
if (write(port, value, count) != count) {
    perror(errtext);
    return(errno);
}
```
```c
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <errno.h>
#include <unistd.h>
#include <locale.h>
#include <time.h>
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>
#define RECIPROCAL(a) (1.0/(a))
#define HALFCYCLE(a) ((a)/2.0)
#define NANOSECONDSPERSECOND 1000000000
int main(int argc, char *argv[]) {
    char *device = "/dev/setparport";
    int port = open(device, O_RDWR);
    char *value;
    int count = 1;
    struct timespec request, remain;
    char errtext[256];
    int i, verbose = 0;
    request.tv_sec = 0;
    request.tv_nsec = 500000; /* 500 microseconds (1 kHz) */
    for (i = 1; i < argc; i++) {
        if (!strcmp(argv[i], "-v"))
            verbose = 1;
        else if (!strcmp(argv[i], "-?")) {
            fprintf(stderr, "Syntax: squarewave \\
                \[<option> \] \[<frequency in Hz>\]\n                \n                Function: Generate square wave at pin #0 of parport (default 1 kHz)\n                \n                Options:   -v show calculated half wave cycle interval (verbose)\n                \n                -? you already figured this one out\n                \n            ");
            exit(1);
        } else {
            char *endptr;
            double freq = strtod(argv[i], &endptr);
            double duration, wholeseconds;
            int endposition = endptr - argv[i];
            if (endposition < strlen(argv[i])) {
                fprintf(stderr, "Could not parse frequency input '\%s' at position \%d (\"c\")\n", argv[i], endposition + 1, argv[i][endposition]);
                exit(1);
            } else {
                duration = HALFCYCLE(RECIPROCAL(freq));
                wholeseconds = (double) (int) duration;
                request.tv_nsec = (int) ((duration - wholeseconds) * NANOSECONDSPERSECOND);
                if (verbose) {
                    setlocale(LC_NUMERIC, "\"\\");
                    printf("Half wave cycle interval is \%d second%s and \%d nanosecond%s\n",
                        request.tv_sec, request.tv_sec == 1 ? "" : "s",
                        request.tv_nsec, request.tv_nsec == 1 ? "" : "s");
                }
            }
        }
    }
    if (port < 0) {
        snprintf(errtext, sizeof(errtext),
            "Could not open device '\%s\', device\n", device);
        perror(errtext);
        return(errno);
    }
    while (1) {
        value = "0";
        if (write(port, value, count) != count) {
            snprintf(errtext, sizeof(errtext),
                "Could not write \%d value%s to device '\%s' at port \%d",
                count, count == 1 ? "" : "s", device, port);
            perror(errtext);
            return(errno);
        }
        clock_nanosleep(CLOCK_MONOTONIC, 0, &request, &remain);
        value = "1";
        write(port, value, 1);
        clock_nanosleep(CLOCK_MONOTONIC, 0, &request, &remain);
    }
    close(port);
}
```

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Calculation

```c
#define RECIPROCAL(a) (1.0/(a))
#define HALFCYCLE(a) ((a)/2.0)
#define NANOSECONDPERSECOND 1000000000

duration = HALFCYCLE(RECIPROCAL(freq));
wholeseconds = (double) ((int) duration);
request.tv_sec = (int) wholeseconds;
request.tv_nsec = (int) ((duration – wholeseconds) * NANOSECONDPERSECOND);
```

Generator

```c
while (1) {
    value = "0";
    clock_nanosleep(CLOCK_MONOTONIC, 0, &request, &remain);
    value = "1";
    write(port, value, 1);
    clock_nanosleep(CLOCK_MONOTONIC, 0, &request, &remain);
}
```
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```
taskset -c 1 squarewave
```

- **RT** C states
- Concurrent process on same CPU core

Overlay: 1 Minute
taskset -c 1 squarewave
taskset -c 1 chrt -f 90 squarewave
taskset -c 1 chrt -f 90 squarewave
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